

## Linear and Integer Optimization

### Exercise Sheet 7

**Exercise 7.1:** Let  $A := \begin{pmatrix} 1 & 0 \\ 0 & 1 \\ s & -1 \end{pmatrix}$  and  $b := \begin{pmatrix} 1 \\ 1 \\ -1 \end{pmatrix}$ .

Use the IDEALIZED ELLIPSOID ALGORITHM with  $R = 2$  to compute a feasible solution in  $P = \{x \in \mathbb{R}^2 \mid Ax \leq b\}$  for  $s = -1$  and for  $s = -2$ . (4 points)

**Exercise 7.2:** Let  $P \subset \mathbb{R}^d$  be a finite non-empty set of points and let  $B$  be a ball containing  $P$ . Show:  $B$  is a ball of minimum radius containing  $P$  if and only if the center of  $B$  lies in  $\text{conv}(P \cap \partial B)$ , where  $\partial B$  is the border of the ball. (5 Points)

**Exercise 7.3:** A semidefinite program is an optimization problem

$$\begin{aligned} \min \quad & C \star X \\ \text{s.t.} \quad & A_i \star X \leq b_i \quad \forall i = 1, \dots, m \\ & X \succeq 0 \\ & X \in \mathbb{R}^{n \times n} \end{aligned}$$

where  $C, A_1, \dots, A_m$  are matrices,  $A \star X := \sum_{1 \leq i, j \leq n} a_{ij} x_{ij}$  and  $X \succeq 0$  means that  $X$  is symmetric and positive semidefinite.

1. Show that the set  $\{X \in \mathbb{R}^{n \times n} \mid X \succeq 0\}$  is a closed cone.
2. Construct a polynomial-time separation oracle for this set. (You may assume that you can compute basic arithmetic operations on real numbers, including square roots, exactly and in constant time.) (3+3 Points)

**Exercise 7.4:** Let  $G$  be a simple undirected graph. Consider the following linear program:

$$\begin{aligned} \min \quad & \sum_{e=\{v,w\} \in E(G)} x_{vw} \\ \text{s.t.} \quad & \sum_{w \in S} x_{vw} \geq \left[ \frac{1}{4}|S|^2 + \frac{1}{2}|S| \right] \quad \text{for } v \in V(G), S \subseteq V(G) \setminus \{v\} \\ & x_{uw} \leq x_{uv} + x_{vw} \quad \text{for } u, v, w \in V(G) \\ & x_{vw} \geq 0 \quad \text{for } v, w \in V(G) \\ & x_{vv} = 0 \quad \text{for } v \in V(G) \end{aligned}$$

- (a) Show that this is a relaxation of the following problem: Find distances  $x_{vw}$  for the nodes of  $G$  such that  $\sum_{e=\{v,w\} \in E(G)} x_{vw}$  is minimized under the condition that there is an ordering  $\{v_1, \dots, v_{|V(G)|}\} = V(G)$  with  $x_{v_i v_j} = |i - j|$  for  $i, j \in \{1, \dots, |V(G)|\}$ .
- (b) Prove that there is a polynomial-time separation oracle for the polyhedron of the feasible solutions of the LP. (3+2 Points)

**Submission deadline:** Tuesday, June 9, 2026, 16:00, via eCampus (in groups of at most 3 students).